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Secure USB bypassing tool

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A B S T R A C T

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As storage capacity increases due to development of flash memory techniques, use of USB memory has increased. As use of USB memory increases, violations of privacy and company confidentiality and technical information leakage occur more often. In this context, use of USB memories that provide security functions to protect the data in them is increasing. Most USB memories are equipped with basic security functions offered by the USB flash drive controller (hereafter called the “USB controller”). However, USB-controller-based security functions have several vulnerabilities. This paper explains how security functions can be bypassed using USB controller commands and presents the design and implementation of a secure USB bypassing tool that bypass the USB security functions.

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1. Introduction

The development of digital technology is making storage devices large in capacity, small in size, and low in unit price. Presently, small storage devices are widely used in daily life, and among these, NAND flash memory devices using the USB (universal serial bus) interface (hereafter called “USB memory”) are the most widely used. USB memories offer a relatively large storage space (64 MB–64 GB) in a physical size smaller than a floppy disk.

However, because of their portability and large capacities, which are their greatest attractions, USB memories are being used in crimes like breaches of company confidentiality and leakage of technical information. In this context, more emphasis is now being given to the importance of USB memories in the course of obtaining evidence during digital forensic investigations (Carvey, 2007).

In an effort to protect against violations of individual privacy that can occur when USB memory is lost, it is becoming more common to protect data using the USB memory security function. The increase in use of USB security is effective in protecting individuals against leakage of private

information, but creates difficulties in obtaining evidence during digital forensic investigations. The reason is that evidence data can be hidden using the security functions of USB memory, and even if the data are known to be hidden, evidence can be obtained only after the suspect’s cooperation has been obtained or a security function bypassing process has been executed. Therefore, if the suspect exhibits an uncooperative attitude when evidence data need to be obtained from secure USB memory, it should be determined whether the security function has been used, and if it has, a process to bypass it is required.

Most USB memories are equipped with a security function using USB controller commands. Because the USB controller security function is easy to use, it is liable to be used for USB-memory-related crimes. However, some USB controller security functions introduce vulnerabilities during the user certification process. Through this process, user certification can be bypassed, and data can be accessed. This paper explains secure USB types, certification methods, and bypassing plans, and presents a design and implementation for a secure USB bypassing tool that can bypass the security function of a secure USB using USB controller commands.

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2. USB security

A secure USB offers a function that blocks an uncertified user from accessing data on the USB and allows use of the data only from a designated PC. Moreover, a secure USB provides various functions like IP tracking that can check the location information of a secure USB user. In general, a secure USB is divided into a non-secure domain with storage space that can be accessed by anyone and a secure domain with storage space that can be accessed only by the designated user.

The security functions of a secure USB include methods to mount hardware for security functions, to certify a user or encode or store data, and to provide security functions using the USB controller commands.

Security hardware is mounted by installing an encoding chip inside and encoding data. In general, AES 128-bit encoding is used, and a strong security function is offered so that if a password is entered incorrectly with more than a certain frequency, the internal security circuit is destroyed or another certification method like fingerprint recognition is used. However, such a system has shortcomings like high price because of the separate hardware mounting.

Security software is a security product which is produced so that it may be used for multiple purposes regardless of USB manufacturer or controller. A user is certified or data are encoded or decoded through internal software commands rather than through the commands built into the USB controller.

The security functions which use USB controller commands, a method of producing security software using the SDK (software development kit) provided by the USB controller manufacturer, are the major topic of analysis in this paper. Data are protected by hiding the partition in the security domain or controlling access as well as by encoding data. Because most USB controllers support security functions and because software based on these functions is distributed for free, users can use them with no difficulty.

3. Secure USB certification vulnerability analysis

USB memory sends data to and receives data from a computer using a device endpoint and a pipe. A device endpoint, a place where data are created or consumed, may be considered as a buffer which stores sent or received data. The pipe serves as a channel to provide and receive data in packet form, in practice by connecting the PC with a device endpoint. The pipe is divided into two parts: the message pipe and the stream pipe. The message pipe sends previously defined commands in packet form, and the stream pipe sends general data (Axelson, 1999).

To detour around the USB security functions, packets transferred between USB memory and the computer are analyzed. Normally the security functions and the relevant packets are sent and received in plain text. In this case, through analysis of sent and received packets, the password can be obtained, and the security function can be cancelled. If the security function and relevant packets are sent and received in code instead of in plain text, obtaining the

Table 1 – Status of Market Shares of USB Controllers in 2007 (Unit: Million Dollars)(2007 USB, 2008).

Order	Manufacturer	Market share	Profit
1	Phison	35.5%	\$32.3
2	Silicon Motion (SMI)	23.2%	\$21.1
3	Sandisk	14.9%	\$13.6
4	Skymedi	9.0%	\$8.2
5	Sony	7.4%	\$6.7
6	AlcorMicro	3.2%	\$2.9
7	Toshiba	3.1%	\$2.8
8	Others	3.7%	\$3.4
	Total	100%	\$91.1

password may be more difficult, but nevertheless, through packet analysis, the security function can be cancelled.

To cancel the security function of a secure USB, which is regarded as a vulnerability to monitoring of USB memory packets, a particular command should be transferred. To transfer commands to a USB or to receive certain information, the DeviceIoControl() API should be used. After receiving the handle of the present USB through the CreateFile() API, data can be sent to and received from the USB using the DeviceIoControl() API.

The commands used in a USB can be divided into common commands used in all USBs and vendor-specific commands that are programmed by the USB controller manufacturer to execute particular functions. The common commands are related to the functions which various USBs commonly execute, like READ and WRITE, and the commands specified by the controller manufacturer are defined in addition for the applicable USB and separately implemented in addition to the common commands. For example, vendor-specific commands used in USB instruments are dependent on the embedded USB controller. Therefore, even if USBs are manufactured by different manufacturers, if the embedded controllers are the same, the same command will be used for the same function, and the security function can be detoured around using the same method.

4. Secure USB certification detouring technique

Among secure USBs using USB controllers, those using three types of controllers, USBest UT163, Skymedi SK6281, and SMI SM321 have certification detouring methods which have

Fig. 1 – USBest UT163 security configuration condition check and password obtaining response message.

06	F9	16	15	03	16	41	58	58	45	4E	00	00	00	53	4B	AXXEN	..SK	
59	20	43	41	4E	44	59	00	00	00	00	00	00	00	31	2E	Y	CANDY1.	
30	30	38	39	38	31	30	33	30	30	30	30	30	30	30	30	00	0089810300000000	=..	
30	30	30	30	30	31	45	42	34	00	02	00	00	00	3D	90	00	0000001EB4..	=..	
00	00	00	00	00	00	TC	45	89	00	08	AA	78	00	00	24	E	.x..\$	
F0	01	00	18	9F	FF	00	00	00	00	00	00	00	00	3D	90	00	00	=..
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
00	00	00	00	00	00	00	00	08	00	00	32	18	04	64	32	00	2..d2	
00	31	32	33	34	00	00	00	00	00	00	00	00	00	00	00	00	1234	
00	31	32	33	34	00	00	00	00	00	00	00	00	00	00	00	00	1234	
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	

Fig. 2 – Skymedi SK6281 security configuration condition checking and password obtaining response message.

already been released (Lee et al., 2008). In addition to these three types of controllers, detouring methods for five other controllers, USBest UT165, Skymedi SK6211, AlcorMicro AU6983, Sandisk Cruzer Contour, and Toshiba, were researched. Table 1 indicates the market shares of several of these USB controllers:

The following section explains the password acquisition and security function detouring methods for each controller.

4.1. USBest UT163

The USBest UT163 controller can perform all of security configuration condition checking, password obtaining, and security certification detouring. Security configuration condition checking and password obtaining can look for the reply message for the command $0 \times F8000000020000001000000$ (12 bytes) after transferring to the USB. This response message has size 0×200 . If the value of the dual offset $0 \times 1E$ (two bytes) is 0×0101 , a security domain exists, otherwise no security domain exists. The password and the password hint may be obtained at offsets 0×28 (8 bytes) and $0 \times E0$ (30 bytes) (Fig. 1).

The command for security certification detouring is issued as $0 \times FC000000000000000000000000000000$ (10 bytes), and the controller cancels the USB security function if it receives the applicable command.

4.2. Skymedi SK6281

The Skymedi SK6281 controller can perform all of security configuration condition checking, password obtaining, and security certification detouring. Security configuration condition checking and password obtaining can look for the response message after transferring the command $0 \times D400000000100000000$ (10 bytes) to the USB. The response message has size 0×200 . If the value of the dual offset $0 \times 3D$ (2 bytes) and that of 0×46 (2 bytes) are different, a security domain exists; if the values are the same, no security domain exists. The password and password hint can be obtained at offset 0×81 (16 bytes) and 0×91 (30 bytes) (Fig. 2).

Fig. 3 – SMI SM321~SM325 security configuration condition checking and password obtaining response message.

[A1] 00 00 00 00 00 00 00 00 |

Fig. 4 – Phison PS2136 security configuration condition checking response message.

When the security certification detouring command is issued as $0 \times D804000000000000$ (10 bytes), the controller cancels the USB security function if it receives the applicable command.

4.3. SMI SM321 ~ SM325

The SMI SM321 ~ SM325 controllers can perform all of security configuration condition checking, password obtaining, and security certification detouring. Security configuration condition checking and password obtaining can look for the response message after transferring the command $0 \times F00500000000000000000000010000000$ (16 bytes) to the USB. The response message has size 0×200 . If the value of dual offset 0×40 (1 byte) is not 0×00 , a security domain exists, and if the value is 0×00 , no security domain exists. The password and password hint can be obtained at offset 0×42 (12 bytes) and 0×50 (16 bytes) (Fig. 3).

When the security certification detouring command is issued as $0 \times F111000000000000000000000000000$ (16 bytes), the controller cancels the USB security function if it receives the applicable command.

4.4. Phison PS2136

The Phison PS2136 controller can perform security configuration condition checking and security certification detouring only. Security configuration condition checking and password obtaining can look for the response message after transferring the command $0 \times 0D0000000800000000000000$ (12 bytes) to the USB. The response message has size 0×08 . If the value of dual offset 0×00 (1 byte) is $0 \times A0$, a security domain exists, and if the value is different, no security domain exists (Fig. 4).

In the Phison PS2136 controller, security certification detouring goes through three stages. First, after transferring the command $0 \times 0605494F464F00000000$ (10 bytes) to the USB, the controller stores the value of offset $0 \times 1FC$ (4 bytes) separately within the response message. Second, the controller replaces offset $0 \times 1FC$ (4 bytes) by 0×00 within the command $0 \times 06060100000000000000000000$ (12 bytes) and the response message from the first process and transfers them simultaneously. Third, the controller enters the value of offset $0 \times 1FC$ (4 bytes) obtained during the first process in reverse order within the [] of $0 \times 0D0000000802[00000000]000000000000$ (16 bytes), the security certification detouring command (Fig. 5).

The Phison PS2136 controller, if the security function detours through certification, loses data from the general domain. Therefore, the data in the general domain should be duplicated in advance before security function certification detouring is performed.

55 53 42 43 10 3B C4 89 10 02 00 00 80 00 10 OD USBC.;
00 00 00 08 02 00 40 3E 00 00 00 00 00 00 00 00 @>

Fig. 5 – Phison PS2136 security function detouring command.

```

43 6F 64 65 4D 61 72 6B 10 39 13 07 01 65 02 08 |CodeMark.9...e..
00 12 89 01 40 8F 08 31 80 10 55 FF 00 00 01 03|....@..1..U.....
FF 80 00 00 FF FF FF FF 38 34 31 30 31 30 00 00|.....841010.
00 00 FF 20 00 A0 0A E0 0A 08 14 00 08 FF FF 00|.....
FC 01 FF FF 06 OA 00 00 00 FF FF FF FF FF FF 01|.....
02 F4 5A 18 58 01 04 04 04 00 00 80 04 00|....Z.X.....
00 80 02 02 1F 00 00 00 55 00 02 40 FF 02 45 B1|....U.@.E...
8B 76 9F AC 60 FF FF 00 01 FF FF FF FF EC E9|....v.....
00 80 02 02 1F 00 00 00 4D 65 6C 52 6F 73 65 20|....Me!Rose...
44 75 61 6C 5F 55 32 30 20 20 20 20 20 20 20|Dual_U20...
30 2E 30 30 FF FF|0.00.....
00 00 00 01 00 02 00 03 00 04 00 05 00 06 00 07|.....
00 08 00 09 00 0A 00 0B FF FF FF FF FF 42 00|....B....
FF FF|.....
62 69 72 74 68 64 61 79 00 00 00 00 00 00 00 00|birthday.....
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00|.....W...
1E 03 30 00 30 00 30 00 30 00 30 00 30 00 30 00|0.0.0.0.0.0.0.

```

Fig. 6 – USBest UT165 security configuration condition checking and password obtaining response message.

4.5. USBest UT163 UT165

The USBest UT163 controller can perform security configuration condition checking and password obtaining only. Security configuration condition checking and password obtaining can look for the response message after transferring the command $0 \times F80000000400000001000000$ (12 bytes) to the USB. The response message has size 0×200 . If the value of dual offset $0 \times 1E$ (2 bytes) is 0×0103 , a security domain exists, and if the value is different, no security domain exists. The password and password hint can be obtained at offset 0×28 (8 bytes) and $0 \times E0(30)$ bytes (Fig. 6).

The USBest UT165 controller transfers the password hash value that uses an independent algorithm along with the security cancellation command to cancel the USB security function. Because it is difficult to know the password hash value in advance, security function detouring is difficult. However, because the password can be found, the security function can be cancelled using the password.

4.6. Skymedi SK6211

The Skymedi SK6211 controller can do all of security configuration condition checking, password obtaining, and security certification detouring. Security configuration condition checking and password obtaining can look for the response message after transferring the command $0 \times D40000000010000000$ (10 bytes) to the USB. The response message has size 0×200 . If the values of dual offset 0×42 (2 bytes) and 0×64 (2 bytes) differ, a security domain exists, and if they have the same value, no security domain exists. The

```

16 03 11 09 E0 3D BD 89 00 02 00 00 80 00 0A CA|.....=.....
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00|00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00|7F.....
30 30 38 39 39 30 30 30 30 30 30 30 30 30 30 31|0089900000AA0401
32 37 30 30 30 30 42 30 45 44 00 01 00 1E F4 00|2700000ED.....
00 00 00 00 00 00 17 7B D6 00 07 78 2A 00 00 1E|...x...{...x*...
F4 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00|.....
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00|00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00|34.....
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00|.....
00 31 32 33 34 00 00 00 00 00 00 00 00 00 00 00|1234.....
00 31 32 33 34 00 00 00 00 00 00 00 00 00 00 00|1234.....
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00|.....

```

Fig. 7 – Skymedi SK6211 security configuration condition checking and password obtaining response message.

```

Send 0x1F bytes to the device
55 53 42 43 B8 72 D4 89 00 02 00 00 00 00 0A 74 |USBC.r.....t
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00|.....
Send 0x200 bytes to the device
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00|.....
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00|.....
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00|.....
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00|.....
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00|.....
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00|.....
77 08 34 32 37 31 28 17 77 33 AA 55 0A 40 6E 67|w.4271(.w3.U.@nq
6F 62 6D 29 63 73 6D AA 55 77 00 00 00 00 00 00|obm)csm.Uw.....
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00|.....

```

Fig. 8 – AlcorMicro AU6983 security function detouring command.

password and password hint may be obtained at offset 0×81 (16 bytes) and 0×91 (30 bytes) (Fig. 7).

When the security certification detouring command is issued as $0 \times F11100000000000000000000000000000$ (16 bytes), the controller cancels the USB security function if it receives the applicable command.

4.7. AlcorMicro AU6983

The AlcorMicro AU6983 controller can perform security configuration condition checking and security certification detouring only. Security configuration condition checking and security certification detouring can look for the response message after transferring the command $0 \times 9C090000020000000000$ (10 bytes) to the USB. The response message has size 0×200 . If the entire response message has the value $0 \times FF$, a security domain does not exist, and if different values are present, a security domain exists.

The AlcorMicro AU6983 controller performs security bypassing in two stages. First, the controller obtains a 512-byte response message to check the security configuration condition. 0×80 (22 bytes) of the response message contain the password hash value used by the independent algorithm. Second, the controller transfers 512 bytes of the response message obtained in the first process simultaneously with $0 \times 74000000000000000000000000000000$ (10 bytes), the security certification detouring command (Fig. 8).

Because the AlcorMicro AU6983 controller uses the password hash value used by the independent algorithm as a password, it is difficult to obtain the password.

4.8. Sandisk Cruzer Contour

The Sandisk Cruzer Contour controller can perform security configuration condition checking and security certification detouring only. Security configuration condition checking can look for the response message after transferring the command $0 \times FFA000000000000000000000000000000$ (16 bytes) to the USB. The response message has size 0×10 . If the values

```

00 60 78 00 00 60 78 00 00 00 00 00 00 00 00 00|`x...`x.....

```

Fig. 9 – Sandisk Cruzer Contour security configuration condition checking response message.

```
Send 0x1F bytes to the device
55 53 42 43 08 90 33 89 00 02 00 00 00 00 00 10 FF | USBC..3.
A2 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00

Send 0x9B bytes to the device
00 60 78 00 C4 10 3F 12 2D 27 67 7C 9D B1 44 CA | `x...?-`g!..D.
E1 39 4A 66 00 00 00 00 00 00 00 00 00 00 00 00 .9jf.....
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
```

Fig. 10 – Sandisk Cruzer Contour security function detouring command transfer screen.

of the dual offsets 0×00 (4 bytes) and 0×04 (4 bytes) are the same, a security domain exists, and if they are different, no security domain exists (Fig. 9).

The Sandisk Cruzer Contour controller processes security bypassing through five stages. First, the controller transfers the command $0 \times FF210000000000000000000000000000$ (16 bytes) to the USB and obtains the value of offset 0×05 (4 bytes) from the response message. Second, the controller substitutes the four-byte value obtained in the first process into the [] of $0 \times [00000000]C4103F122D27677C9DB144CAE1394A66$ (16 bytes) and saves it separately. Third, the controller transfers $0 \times FFA200000000000000000000000000000$ (16 bytes), the security bypassing command, and the value saved in the second process simultaneously. Fourth, the controller transfers $0 \times FF210000000000000000000000000000$ (16 bytes), the last checking command (Fig. 10).

Because the Sandisk Cruzer Contour controller uses the password hash value that used by the independent algorithm as a hash, obtaining the password is difficult. In addition, if certification is detoured around using the security bypassing method previously described, the password is changed to an arbitrary password.

4.9. Toshiba

The Toshiba controller can perform security configuration condition checking and security bypassing only. Security configuration condition checking can look for the response message after transferring the command $0 \times 1A003F000C0000000000$ (10 bytes) to the USB. The response message has size 0×200 . If the value of dual offset 0×25 (1 byte) is 0×00 , a security domain exists, and if it is 0×01 , no security domain exists (Fig. 11).

To detour around the security certification of the Toshiba controller, the current password should be modified to an arbitrary value. The Toshiba controller changes the password to the 20-byte hash value using an independent algorithm. In this paper, the certification detouring process will be explained using the value 0 × 8BDFA5B53A0-F5485A8933C7961891F7D1894997B (20 bytes) which is

65 44 09 30 00 00 00 3B AF C0 00 3B AF C0 00 3B eD.0...;.;.;.;.
AF C0 00 3B AF C0 00 00 00 00 00 22 90 00 00 22;"."
90 00 00 19 18 00 00 3B AF C0 00 00 00 00 00 00 00,.....
00 00 00 00 00 00 FF
FF FF FF FF FF FF FF FF FF FF FF FF FF FF FF FF FF FF
00 FF
00 54 5F 53 48 49 42 41 20 54 72 61 6E 73 4D 65 TOSHIBA TransMe

Fig. 11 – Toshiba security configuration condition checking response message.

```
Send 0x1F bytes to the device
55 53 42 43 90 E4 C1 89 00 02 00 00 00 00 0A FF USBC...
57 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 W.......
```



```
Send 0x200 bytes to the device
8B DF 45 B5 3A 0F 54 85 A8 93 3C 79 61 89 1F 7D ...;T...<ya...>
18 94 99 7B 00 00 00 00 00 00 00 00 00 00 00 00 00 ...{.....
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 .....
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 .....
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 .....
```

Fig. 12 – Toshiba security function detouring command transfer screen.

equivalent to '1234'. The certification detouring process has two stages. First, the process simultaneously transfers the password-modifying command $0 \times FF570000000000000000$ (10 bytes) and the value $0 \times 8BDFA5B53A0F5485A8933C7961891F7D1894997B$ (20 bytes), the hash value equivalent to '1234'. This serves to change the current password to '1234'. Second, the process transfers simultaneously the security bypassing command $0 \times FF540000000000000000$ (10 bytes) and the hash value equivalent to '1234' (Fig. 12).

Whether the password is obtained, and whether the security certification can be detoured around (Table 2).

5. Design and implementation of secure usb certification detouring tools

In this research, a set of secure USB certification detouring tools was designed and implemented using the vulnerabilities of a secure USB and USB controller commands. A secure USB certification detouring tool can detour around the security function of a USB controller. Some of its specific functions include:

Table 2 – Whether the security function can be detoured around by various USB controllers.

Controller name	Security configuration condition checking	Whether password is obtained	Whether security certification can be detoured around
USBest UT163	O	O	O
Skymedi SK6281	O	O	O
SMI SM321 ~ SM325	O	O	O
Phison PS2136	O	X	O
USBest UT165	O	O	X
Skymedi SK6211	O	O	O
AlcorMicro AU6983	O	X	O
Sandisk Cruzer Contour	O	X	O
Toshiba	O	X	O

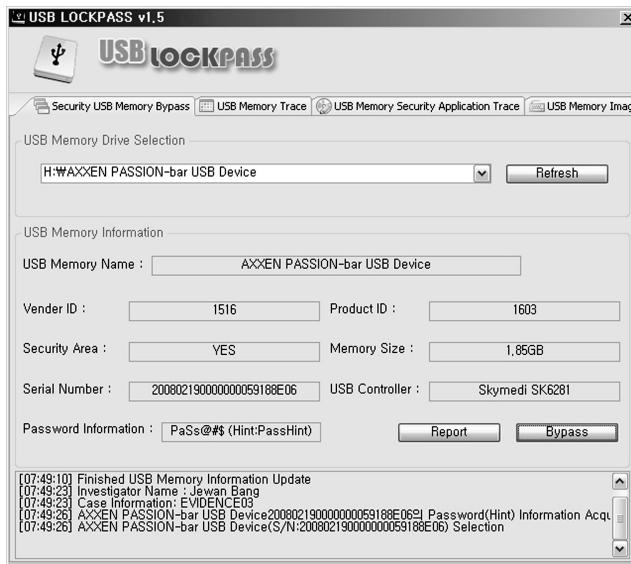


Fig. 13 – Certification detouring screen of the secure USB bypassing tool.

- Secure USB security function canceling function
 - Secure USB device information analysis function
 - Secure USB password analysis function
 - USB memory write protection function
 - Function to analyze the information in the USB memory mounted in the intended system
 - Function to analyze the traces of using the security canceling tool of the secure USB
 - Disk image creation function
 - Report composition function.

If the security function is active in an USB which has been obtained as evidence in a digital forensic investigation, the USB certification detouring tool can obtain data from this USB by detouring around the secure USB certification. In addition,

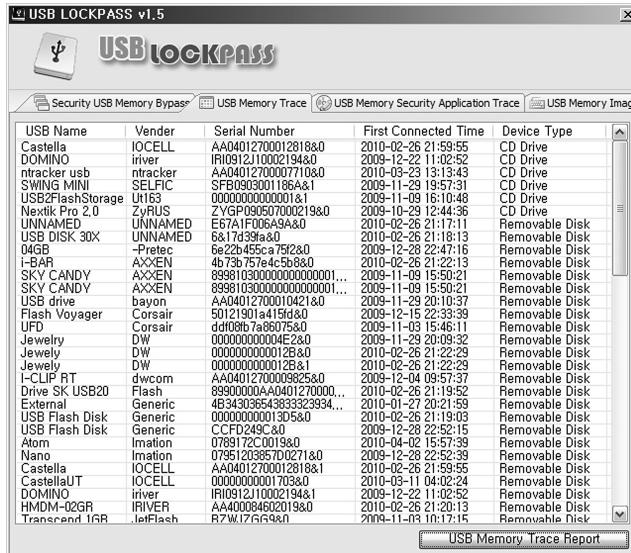


Fig. 14 – Trace screen of the secure USB certification bypassing tool.

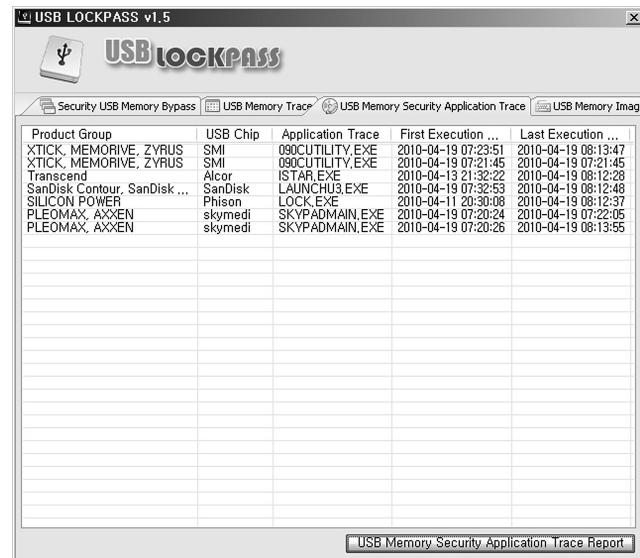


Fig. 15 – Trace analysis function of the secure USB bypassing tool’s security cancellation tool.

this tool can create an image of the data obtained and compose a report on the data. In addition, the USB certification detouring tool can analyze secure USB device information and passwords and the information in the USB memory mounted in the target system using information saved in the registry(Carvey and Altheide, 2005). It also provides a USB memory write protection function to prevent corruption of the target USB memory during analysis. The USB memory's write protection function can activate this function by changing the configuration of the registry. If WriteProtectValue of "HKEY_LOCAL_MACHINE\SYSTEM\CurrentControlSet\Control\Storage DevicePolicies" is set to one, the write protection function is activated (Figs. 13 and 14).

Fig. 17 shows the structure of the USB bypassing tool. The tool is composed of three parts: the USB device recognition

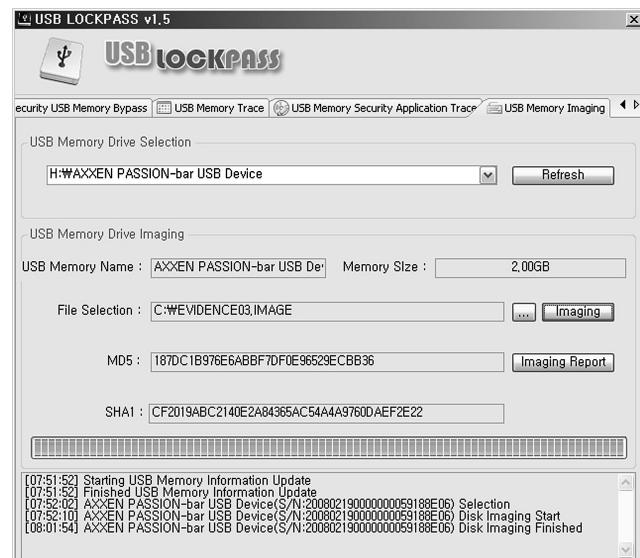


Fig. 16 – USB imaging screen of secure USB bypassing tool.

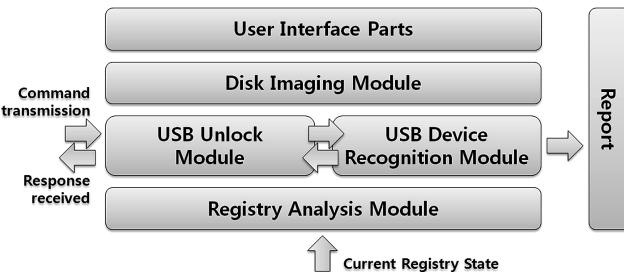


Fig. 17 – Structure of the secure USB bypassing tool.

module, the USB unlock module, and the registry analysis module.

The USB device recognition module analyzes and outputs basic information by recognizing the USB mounted in the target drive. The USB unlock module cancels the USB security functions based on the information transferred from the USB device recognition module and executes other operations such as password extraction. The registry analysis module reads the registry information of the target system and analyzes the information on any USB memory mounted in the system (Figs. 15–17).

6. Conclusions

In this paper, methods of bypassing USB security functions using USB controller commands have been explained, and the design and implementation of a secure USB bypassing tool that bypasses USB security functions has been described. As mentioned in the introduction, the use of USB memory has increased as USB memories have evolved toward larger capacities and small sizes. That means a greater importance for USB memory in the course of obtaining evidence in digital forensic investigations. Furthermore, if USB security functions are active in a USB obtained as evidence, the process of detouring around the security functions of a secure USB becomes more important to obtain evidence.

Therefore, in addition to the detouring method for USB controller security functions proposed in this paper, the security functions of new USB controllers to be launched in the future should be researched.

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